

Tutorial 8

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¹based on material provided by Diane Horton, Eric Joanis, Alfredo Gabaldon

1 Lists in Prolog

count(L,E,N) the list L contains N copies of E.
Precondition: L and E are instantiated.

```
count([],_,0).  
count([E|Rest],E,N) :- count(Rest,E,N1), N is N1 + 1.  
count([X|Rest],E,N) :- \+(X=E), count(Rest, E, N).
```

element_at(X,L,K) X is the kth element in list L; (start with 1)
Precondition: K is instantiated.

```
element_at(X,[X|_],1).  
element_at(X,[_|L],K) :- K > 1, K1 is K - 1, element_at(X,L,K1).
```

Another solution: K does not have to be instantiated.

```
element_at(X,[X|_],1).  
element_at(X,[_|L],K) :- element_at(X,L,K1), K is K1 + 1.
```

Note: The 1st solution never goes into infinite loop,
 but it fails if K is a variable.

`delete(E,L1,L2)` list L2 is list L1 with exactly one instance of E "deleted".

```
delete(E,[E|Rest],Rest).
```

```
delete(E,[X|Rest],[X|Rest2]) :- delete(E,Rest,Rest2).
```

Notice: we don't have $X \neq E$ in the recursive case.

That's because the query

```
?- delete(1,[1,2,1],L).
```

should say

```
L = [2,1];
```

```
L = [1,2];
```

```
no
```

A side-effect of this is that the query

```
?- delete(1,[1,1],L).
```

will also say

```
L = [1];
```

```
L = [1];
```

```
no
```

(deleting the second copy of 1).

reverse(L1,L2) L2 is the list L1 in reverse order
Precondition: L1 is instantiated.

```
reverse([],[]).  
reverse([A|X],Z) :- reverse(X,Y), append(Y,[A],Z).
```

Notice the use of 'append' here.
This solution is 'slow' - quadratic time.

Draw a Prolog search tree for:

```
reverse([1, 2, 3], L).
```

```
permutation(L1,L2)    L2 is some permutation of L1,  
                    i.e., it has the same elements, but in any order.  
Precondition:       L1 is instantiated.
```

```
permutation([],[]).  
permutation(L1,[X|Rest2]) :- delete(X,L1,Rest1),    % defined earlier  
                             permutation(Rest1,Rest2).
```

Another Solution:

```
permutation([],[]).  
permutation(L1,[X|R2]) :- append(P1,[X|P2],L1),  
                          append(P1,P2,R1),  
                          permutation(R1,R2).
```

Again, notice how 'append' is used here.

2 Additional Material

`reverse(L1,L2)` L2 is the list L1 in reverse order
Precondition: L1 is instantiated.

Solution 2:

```
reverse(X,Z) :- reverse_acc(X,[],Z).
```

```
reverse_acc([],Y,Y).
```

```
reverse_acc([A|X],Y,Z) :- reverse_acc(X,[A|Y],Z).
```

Y is called an ‘‘accumulator’’ variable.
This solution is ‘‘fast’’ -- linear time.

Draw a search tree for :

```
reverse([1,2,3],L)
```

Draw a search tree for :

`reverse(L, [1,2,3])`

We get infinite recursion after giving out the answer.

Thus, we need a precondition:

the first list is instantiated to a list of known length,
i.e., the tail is [], and not an uninstantiated variable.

The following implementation works with any input:

```
reverse(L,R) :- knownlength(L), reverse_acc(L,[],R).  
reverse(L,R) :- \+knownlength(L), reverse_acc(R,[],L).
```

```
knownlength(L) :- \+var(L),  
  (L=[]; L=[_|R], knownlength(R)).
```

```
reverse_acc([],Y,Y).  
reverse_acc([A|X],Y,Z) :- reverse_acc(X,[A|Y],Z).
```

Draw a search tree for:

```
reverse(L,[1,2,3])
```

to see that it now works.