

Lecture 17: Formal Modeling Methods

Formal Modeling Techniques

Definition of FM

Why use FM?

Program Specification vs. Requirements Modeling

Example Formal Methods:

RSML

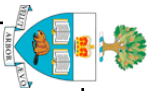
SCR

RML

Telos

Albert II

Tips on formal modeling



What are Formal Methods?

Broad View (Leveson)

application of discrete mathematics to software engineering
involves modeling and analysis
with an underlying mathematically-precise notation

Narrow View (Wing)

Use of a formal language
a set of strings over some well-defined alphabet, with rules for distinguishing which strings belong to the language

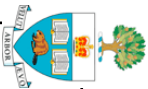
Formal reasoning about formulae in the language

E.g. formal proofs: use axioms and proof rules to demonstrate that some formula is in the language

For requirements modeling...

A notation is formal if:

- ...it comes with a formal set of rules which define its syntax and semantics.
- ...the rules can be used to analyse expressions to determine if they are syntactically well-formed or to prove properties about them.



Formal Methods in Software Engineering

What to formalize?

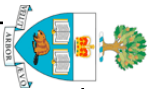
models of requirements knowledge (so we can reason about them)
specifications of requirements (so we can document them precisely)
Specifications of program design (so we can verify correctness)

Why formalize?

Removes ambiguity and improves precision
To verify that the requirements have been met
To reason about the requirements/designs
Properties can be checked automatically
Test for consistency, explore consequences, etc.
To animate/execute specifications
Helps with visualization and validation
...because we have to formalize eventually anyway
Need to bridge from the informal world to a formal machine domain

Why people don't formalize!

Formal Methods tend to be lower level than other techniques
They include too much detail
Formal Methods concentrate on consistent, correct models
...most of the time your models are inconsistent, incorrect, incomplete...
People get confused about which tools are appropriate:
specification of program behaviour vs. modeling of requirements
formal methods advocates get too attached to one tool!
Formal methods require more effort
...and the payoff is deferred



Consistency analysis and typechecking

"Is the formal model well-formed?"

Assumes "well-formedness" of the model corresponds to something useful...

Validation:

Animate the model on small examples

Formal challenges:

"if the model is correct then the following property should hold..."

'what if' questions:

reasoning about the consequences of particular requirements;
reasoning about the effect of possible changes

Predicting behavior

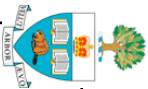
State exploration (E.g. through model checking)

Checking application properties:

"will the system ever do the following..."

Verifying design refinement

"does the design meet the requirements?"



Three traditions ...

Formal Specification Languages

Grew out of work on program verification

Spawned many general purpose specification languages

Good for specifying the behaviour of program units

Key technologies: Type checking, Theorem proving

Applicable to program design

- closely tied to program semantics

Examples: Larch, Z, VDM, ...

Reactive System Modelling

Formalizes dynamic models of system behaviour

Good for reactive systems (e.g. real-time, embedded control systems)

can reason about safety, liveness, performance(?)

Key technologies: Consistency checking, Model checking

Applicable to Requirements

- Languages developed specifically for RE

Examples: Statecharts, RSML, *Parnas-tables, SCR, ...*

Formal Conceptual Modelling

For capturing real-world knowledge in RE

Focuses on modelling domain entities, activities, agents, assertions, goals,...

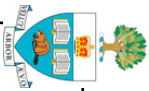
use first order predicate logic as the underlying formalism

Key technologies: inference engines, default reasoning, KBS-shells

Applicable to Requirements

- Capture key requirements concepts

Examples: Reqts Apprentice, RML, Telos, Albert II, ...



(1) Formal *Specification* Languages

Three basic flavours:

Operational - specification is executable abstraction of the implementation
good for rapid prototyping

e.g., Lisp, Prolog, Smalltalk

State-based - views a program as a (large) data structures whose state can be altered by procedure calls...

... using pre/post-conditions to specify the effect of procedures

e.g., VDM, Z

Algebraic - views a program as a set of abstract data structures with a set of operations...

... operations are defined declaratively by giving a set of axioms

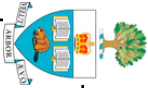
e.g., Larch, CLEAR, OBJ

Developed for specifying *programs*

Programs are formal, man-made objects

... and can be modeled precisely in terms of input-output behaviour

These languages are NOT appropriate for requirements modeling
requirements specification \neq program specification



(2) Reactive System *Modelling*

Modeling how a system should behave

General approach:

Model the environment as a state machine

Model the system as a state machine

Model safety, liveness properties of the machine as temporal logic assertions

Check whether the properties hold of the system interacting with its environment

Examples:

Statecharts

Harel's notation for modeling large systems

Adds parallelism, decomposition and conditional transitions to STDs

RSML

Heimdahl & Leveson's Requirements State Machine Language

Adds tabular specification of complex conditions to Statecharts

A7e approach

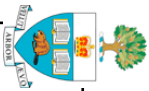
Major project led by Parnas to formalize A7e aircraft requirements spec

Uses tables to specify transition relations & outputs

SCR

Heitmeyer et. al. "Software Cost Reduction"

Extends the A7e approach to include dictionaries & support tables



(3) Formal Conceptual Modelling

General approach

model the world beyond software functions
build models of humans' knowledge/beliefs about the world
draws on techniques from AI and Knowledge Representation
make use of abstraction & refinement as structuring primitives

Examples:

RML - Requirements Modeling Language

Developed by Greenspan & Mylopoulos in mid-1980s
First major attempt to use knowledge representation techniques in RE
Object oriented language, with classes for activities, entities and assertions
Uses First Order Predicate Language as an underlying reasoning engine

Telos

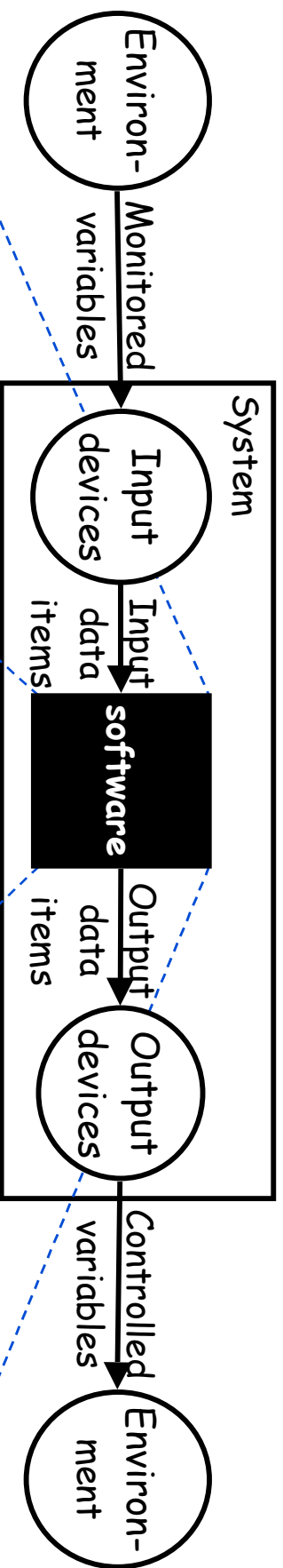
Extends RML by creating a fully extensible ontology
meta-level classes define the ontology (the basic set is built in)

Albert II

developed by Dubois & du Bois in the mid-1990s
Models a set of interacting **agents** that perform **actions** that change their **state**
uses an object-oriented real-time temporal logic for reasoning

Example: SCR

Four Variable Model:



Dictionary:

[illegible]

Type	Base Value	Units	Warming Level	Humidity Level	Mid-high Temperature

Constant Type Value	Unit	Storage	High Temperature
1	1	1	1
2	2	2	2
3	3	3	3
4	4	4	4
5	5	5	5
6	6	6	6
7	7	7	7
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Tables:

[illegible]

Event Tables

a/s/o:

Assertions, Scenarios,

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Monitored/Controlled

Variables

[illegible]

Mode Transition Tables

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Event Tables

Event Tables

Assertions,

a/s/o:

Types

Type	Base Value	Units	Warming Level	Humidity Level	Mid-high Temperature

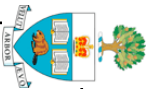
Constants

Constant Type Value	Unit	Storage	High Temperature
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Condition Tables

[illegible]

SCR Specification



SCR basics

Source: Adapted from Heimeyer et. al. 1996.

Modes and Mode classes

A mode class is a finite state machine, with states called *system modes*

Transitions in each mode class are triggered by events

Complex systems are described using a number of mode classes operating in parallel

System State

A (system) state is defined as:

the system is in exactly one mode from each mode class...

...and each variable has a unique value

Events

An event occurs when any system entity changes value

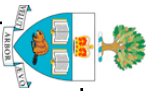
An input event occurs when an input variable changes value

Single input assumption - only one input event can occur at once

Notation: @T(c) means "c changed from false to true"

A conditioned event is an event with a predicate

@T(c) WHEN d means: "c became true when c was false and d was true"



SCR Tables

Source: Adapted from Heimeyer et. al. 1996.

Mode Class Tables

Define the set of *modes* (states) that the software can be in.

A complex system will have many different modes classes

Each mode class has a mode table showing the conditions that cause transitions between modes

A mode table defines a *partial function* from modes and events to modes

Event Tables

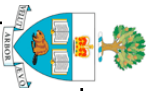
An event table defines how a term or controlled variable changes in response to input events

Defines a *partial function* from modes and events to variable values

Condition Tables

A condition table defines the value of a term or controlled variable under every possible condition

Defines a *total function* from modes and conditions to variable values

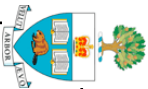


Example: Temp Control System

Source: Adapted from Heitmeyer et. al. 1996.

Mode transition table:

Current Mode	Powered on	Too Cold	Temp OK	Too Hot	New Mode
Off	@T	-	t	-	Inactive
	@T	t	-	-	Heat
	@T	-	-	t	AC
Inactive	@F	-	-	-	Off
	-	@T	-	-	Heat
	-	-	-	@T	AC
Heat	@F	-	-	-	Off
	-	-	@T	-	Inactive
AC	@F	-	-	-	Off
	-	-	@T	-	Inactive



Failure modes

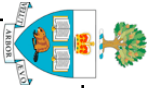
Source: Adapted from Heimeyer et. al. 1996.

Mode transition table:

Current Mode	Powered on	Cold Heater	Too Cold	Warm AC	Too Hot	New Mode
NoFailure	t	@T	t	-	-	HeatFailure
HeatFailure	t	-	-	@T	t	ACFailure
ACFailure	t	@F	t	-	-	NoFailure
ACFailure	t	-	-	@F	t	NoFailure

Event table:

Modes		
NoFailure	@T(INMODE)	never
ACFailure, HeatFailure	never	@T(INMODE)
Warning light =	Off	On



Using Formal Methods

Selective use of Formal Methods

Amount of formality can vary

Need not build complete formal models

Apply to the most critical pieces

Apply where existing analysis techniques are weak

Need not formally analyze every system property

E.g. check safety properties only

Need not apply FM in every phase of development

E.g. use for modeling requirements, but don't formalize the system design

Can choose what level of abstraction (amount of detail) to model

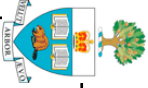
Lightweight Formal Methods

Have become popular as a means of getting the technology transferred

Two approaches

Lightweight *use of* FMs - selectively apply FMs for partial modeling

Lightweight FMs - new methods that allow unevaluated predicates



References

van Vliet, H. "Software Engineering: Principles and Practice (2nd Edition)" Wiley, 1999.

van Vliet gives a good introduction to formal methods in chapter 15. In particular, sections 15.1 and 15.5 are worth reading, to give a feel for the current state of the art, and the problems that hinder the use of formal methods in practice. van Vliet describes a completely different set of formal modeling techniques from those covered in this lecture – he concentrates on methods that can be used for program design models, rather than requirements models.

Heitmeyer, C. L., Jeffords, R. D., & Labaw, B. G. (1996). Automated Consistency Checking of Requirements Specifications. ACM Transactions on Software Engineering and Methodology, 5(3), 231–261.

Describes SCR in detail.